Abstract

This document specifies an architectural framework for the application of Multiprotocol Label Switching (MPLS) to the construction of packet-switched transport networks. It describes a common set of protocol functions - the MPLS Transport Profile (MPLS-TP) - that supports the operational models and capabilities typical of such networks, including signaled or explicitly provisioned bi-directional connection-oriented paths, protection and restoration mechanisms, comprehensive Operations, Administration and Maintenance (OAM) functions, and network operation in the absence of a dynamic control plane or IP forwarding support. Some of these functions are defined in existing MPLS specifications, while others require extensions to existing specifications to meet the requirements of the MPLS-TP.

This document defines the subset of the MPLS-TP applicable in general and to point-to-point paths. The remaining subset, applicable specifically to point-to-multipoint paths, are out of scope of this document.

This document is a product of a joint Internet Engineering Task Force (IETF) / International Telecommunications Union Telecommunications Standardization Sector (ITU-T) effort to include an MPLS Transport Profile within the IETF MPLS and PWE3 architectures to support the capabilities and functionalities of a packet transport network as defined by the ITU-T.

Status of This Memo

This Internet-Draft is submitted to IETF in full conformance with the provisions of BCP 78 and BCP 79.
Internet-Drafts are working documents of the Internet Engineering Task Force (IETF), its areas, and its working groups. Note that other groups may also distribute working documents as Internet-Drafts.

Internet-Drafts are draft documents valid for a maximum of six months and may be updated, replaced, or obsoleted by other documents at any time. It is inappropriate to use Internet-Drafts as reference material or to cite them other than as "work in progress."

The list of current Internet-Drafts can be accessed at http://www.ietf.org/ietf/1id-abstracts.txt.

The list of Internet-Draft Shadow Directories can be accessed at http://www.ietf.org/shadow.html.

This Internet-Draft will expire on August 8, 2010.

Copyright Notice

Copyright (c) 2010 IETF Trust and the persons identified as the document authors. All rights reserved.

This document is subject to BCP 78 and the IETF Trust’s Legal Provisions Relating to IETF Documents (http://trustee.ietf.org/license-info) in effect on the date of publication of this document. Please review these documents carefully, as they describe your rights and restrictions with respect to this document. Code Components extracted from this document must include Simplified BSD License text as described in Section 4.e of the Trust Legal Provisions and are provided without warranty as described in the BSD License.

Table of Contents

1. Introduction ................................................. 4
   1.1. Motivation and Background ................................. 4
   1.2. Scope .................................................. 5
   1.3. Terminology ............................................. 6
   1.3.1. Transport Network ................................... 6
   1.3.2. MPLS Transport Profile ................................. 7
   1.3.3. MPLS-TP Section ..................................... 7
   1.3.4. MPLS-TP Label Switched Path ........................... 7
   1.3.5. MPLS-TP Label Switching Router (LSR) and Label Edge Router (LER) ............................. 7
   1.3.6. Customer Edge (CE) .................................. 8
   1.3.7. Edge-to-Edge LSP ...................................... 8
   1.3.8. Service LSP .......................................... 8
1. Introduction

1.1. Motivation and Background

This document describes an architectural framework for the application of MPLS to the construction of packet-switched transport networks. It specifies the common set of protocol functions that meet the requirements in [RFC5654], and that together constitute the MPLS Transport Profile (MPLS-TP) for point-to-point paths. The remaining MPLS-TP functions, applicable specifically to point-to-multipoint paths, are out of scope of this document.

Historically the optical transport infrastructure - Synchronous Optical Network/Synchronous Digital Hierarchy (SONET/SDH) and Optical Transport Network (OTN) - has provided carriers with a high benchmark for reliability and operational simplicity. To achieve this, transport technologies have been designed with specific characteristics:

- Strictly connection-oriented connectivity, which may be long-lived and may be provisioned manually (i.e. configuration of the node via a command line interface) or by network management.
- A high level of availability.
- Quality of service.
- Extensive OAM capabilities.

Carriers wish to evolve such transport networks to take advantage of the flexibility and cost benefits of packet switching technology and to support packet based services more efficiently. While MPLS is a maturing packet technology that already plays an important role in transport networks and services, not all MPLS capabilities and mechanisms are needed in or consistent with the transport network operational model. There are also transport technology characteristics that are not currently reflected in MPLS.

There are thus two objectives for MPLS-TP:

1. To enable MPLS to be deployed in a transport network and operated in a similar manner to existing transport technologies.

2. To enable MPLS to support packet transport services with a similar degree of predictability to that found in existing transport networks.

In order to achieve these objectives, there is a need to define a
common set of MPLS protocol functions - an MPLS Transport Profile - for the use of MPLS in transport networks and applications. Some of the necessary functions are provided by existing MPLS specifications, while others require additions to the MPLS tool-set. Such additions should, wherever possible, be applicable to MPLS networks in general as well as those that conform strictly to the transport network model.

This document is a product of a joint Internet Engineering Task Force (IETF) / International Telecommunications Union Telecommunications Standardization Sector (ITU-T) effort to include an MPLS Transport Profile within the IETF MPLS and PWE3 architectures to support the capabilities and functionalities of a packet transport network as defined by the ITU-T.

1.2. Scope

This document describes an architectural framework for the application of MPLS to the construction of packet-switched transport networks. It specifies the common set of protocol functions that meet the requirements in [RFC5654], and that together constitute the MPLS Transport Profile (MPLS-TP) for point-to-point MPLS-TP transport paths. The remaining MPLS-TP functions, applicable specifically to point-to-multipoint transport paths, are out of scope of this document.
1.3. Terminology

<table>
<thead>
<tr>
<th>Term</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>LSP</td>
<td>Label Switched Path</td>
</tr>
<tr>
<td>MPLS-TP</td>
<td>MPLS Transport Profile</td>
</tr>
<tr>
<td>SDH</td>
<td>Synchronous Digital Hierarchy</td>
</tr>
<tr>
<td>ATM</td>
<td>Asynchronous Transfer Mode</td>
</tr>
<tr>
<td>OTN</td>
<td>Optical Transport Network</td>
</tr>
<tr>
<td>cl-ps</td>
<td>Connectionless - Packet Switched</td>
</tr>
<tr>
<td>co-cs</td>
<td>Connection Oriented - Circuit Switched</td>
</tr>
<tr>
<td>co-ps</td>
<td>Connection Oriented - Packet Switched</td>
</tr>
<tr>
<td>OAM</td>
<td>Operations, Administration and Maintenance</td>
</tr>
<tr>
<td>G-ACh</td>
<td>Generic Associated Channel</td>
</tr>
<tr>
<td>GAL</td>
<td>G-ACh Label</td>
</tr>
<tr>
<td>MEP</td>
<td>Maintenance End Point</td>
</tr>
<tr>
<td>MIP</td>
<td>Maintenance Intermediate Point</td>
</tr>
<tr>
<td>APS</td>
<td>Automatic Protection Switching</td>
</tr>
<tr>
<td>SCC</td>
<td>Signaling Communication Channel</td>
</tr>
<tr>
<td>MCC</td>
<td>Management Communication Channel</td>
</tr>
<tr>
<td>EMF</td>
<td>Equipment Management Function</td>
</tr>
<tr>
<td>FM</td>
<td>Fault Management</td>
</tr>
<tr>
<td>CM</td>
<td>Configuration Management</td>
</tr>
<tr>
<td>PM</td>
<td>Performance Management</td>
</tr>
<tr>
<td>LSR</td>
<td>Label Switching Router</td>
</tr>
<tr>
<td>MPLS-TP PE</td>
<td>MPLS-TP Provider Edge LSR</td>
</tr>
<tr>
<td>MPLS-TP P</td>
<td>MPLS-TP Provider LSR</td>
</tr>
<tr>
<td>PW</td>
<td>Pseudowire</td>
</tr>
<tr>
<td>AC</td>
<td>Attachment Circuit</td>
</tr>
<tr>
<td>Adaptation</td>
<td>The mapping of client information into a format suitable for transport by</td>
</tr>
<tr>
<td></td>
<td>the server layer</td>
</tr>
<tr>
<td>Native</td>
<td>The traffic belonging to the client of the MPLS-TP network</td>
</tr>
<tr>
<td>Service</td>
<td></td>
</tr>
<tr>
<td>T-PE</td>
<td>PW Terminating Provider Edge</td>
</tr>
<tr>
<td>S-PE</td>
<td>PW Switching provider Edge</td>
</tr>
</tbody>
</table>

1.3.1. Transport Network

A Transport Network provides transparent transmission of client user plane traffic between attached client devices by establishing and maintaining point-to-point or point-to-multipoint connections between such devices. The architecture of networks supporting point to multipoint connections is out of scope of this document. A Transport Network is independent of any higher-layer network that may exist between clients, except to the extent required to supply this transmission service. In addition to client traffic, a Transport Network may carry traffic to facilitate its own operation, such as that required to support connection control, network management, and
Operations, Administration and Maintenance (OAM) functions.

See also the definition of Packet Transport Service in Section 3.1.

1.3.2.  MPLS Transport Profile

The MPLS Transport Profile (MPLS-TP) is the subset of MPLS functions that meet the requirements in [RFC5654]. Note that MPLS is defined to include any present and future MPLS capability specified by the IETF, including those capabilities specifically added to support transport network requirements [RFC5654].

1.3.3.  MPLS-TP Section

An MPLS-TP Section is defined in Section 1.2.2 of [RFC5654].

1.3.4.  MPLS-TP Label Switched Path

An MPLS-TP Label Switched Path (MPLS-TP LSP) is an LSP that uses a subset of the capabilities of an MPLS LSP in order to meet the requirements of an MPLS transport network as set out in [RFC5654]. The characteristics of an MPLS-TP LSP are primarily that it:

1. Uses a subset of the MPLS OAM tools defined as described in [I-D.ietf-mpls-tp-oam-framework].
2. Supports 1+1, 1:1, and 1:N protection functions.
3. Is traffic engineered.
4. May be established and maintained via the management plane, or using GMPLS protocols when a control plane is used.
5. Is either point-to-point or point-to-multipoint. Multipoint to point and multipoint to multipoint LSPs are not permitted.

Note that an MPLS LSP is defined to include any present and future MPLS capability, including those specifically added to support the transport network requirements.

1.3.5.  MPLS-TP Label Switching Router (LSR) and Label Edge Router (LER)

An MPLS-TP Label Switching Router (LSR) is either an MPLS-TP Provider Edge (PE) router or an MPLS-TP Provider (P) router for a given LSP, as defined below. The terms MPLS-TP PE router and MPLS-TP P router describe logical functions; a specific node may undertake only one of these roles on a given LSP.
Note that the use of the term "router" in this context is historic and neither requires nor precludes the ability to perform IP forwarding.

1.3.5.1. MPLS-TP Provider Edge (PE) Router

An MPLS-TP Provider Edge (PE) router is an MPLS-TP LSR that adapts client traffic and encapsulates it to be transported over an MPLS-TP LSP. Encapsulation may be as simple as pushing a label, or it may require the use of a pseudowire. An MPLS-TP PE exists at the interface between a pair of layer networks. For an MS-PW, an MPLS-TP PE may be either an S-PE or a T-PE, as defined in [RFC5659].

1.3.5.2. MPLS-TP Provider (P) Router

An MPLS-TP Provider router is an MPLS-TP LSR that does not provide MPLS-TP PE functionality for a given LSP. An MPLS-TP P router switches LSPs which carry client traffic, but does not adapt client traffic and encapsulate it to be carried over an MPLS-TP LSP.

1.3.5.3. Label Edge Router (LER)

An LSR that exists at the endpoints of an LSP and therefore pushes or pops a label, i.e. does not perform a label swap on the particular LSP under consideration.

1.3.6. Customer Edge (CE)

A Customer Edge (CE) is the client function sourcing or sinking native service traffic to or from the MPLS-TP network. CEs on either side of the MPLS-TP network are peers and view the MPLS-TP network as a single point-to-point or point-to-multipoint link.

1.3.7. Edge-to-Edge LSP

An Edge-to-Edge LSP is an LSP between a pair of PEs that may transit zero or more provider LSRs.

1.3.8. Service LSP

A service LSP is an LSP that carries a single client service.

1.3.9. Layer Network

A layer network is defined in [G.805] and described in [RFC5654].
1.3.10. Additional Definitions and Terminology

Detailed definitions and additional terminology may be found in [RFC5654].

1.4. Applicability

MPLS-TP can be used to construct packet transport networks and is therefore applicable in any packet transport network context. It is also applicable to subsets of a packet network where the transport network operational model is deemed attractive. The following are examples of MPLS-TP applicability models:

1. MPLS-TP provided by a network that only supports MPLS-TP LSPs and PWs (i.e. Only MPLS-TP LSPs and PWs exist between the PEs or LSRs), acting as a server for other layer 1, layer 2 and layer 3 networks (Figure 1).

2. MPLS-TP provided by a network that also supports non-MPLS-TP LSPs and PWs (i.e. both LSPs and PWs that conform to the transport profile and those that do not, exist between the PEs), acting as a server for other layer 1, layer 2 and layer 3 networks (Figure 2).

3. MPLS-TP as a server layer for client layer traffic of IP or MPLS networks which do not use functions of the MPLS transport profile. For MPLS traffic, the MPLS-TP server layer network uses PW switching [RFC5659] or LSP stitching [RFC5150] at the PE that terminates the MPLS-TP server layer (Figure 3).

These models are not mutually exclusive.
MPLS-TP LSP, provided by a network that only supports MPLS-TP, acting as a server for other layer 1, layer 2 and layer 3 networks.

```plaintext
<-- L1/2/3 -->|<-- MPLS-TP-->|<-- L1/2/3 -->|

Only

MPLS-TP
++++   LSP   ++++
++++  Client  +------ Client  ++++
|CE1=Traffic=|PE2=---------|PE3=Traffic=|CE1|
++++  +------ +------ +------ +------ +------ +++

Example
a)  [Ethernet] [Ethernet] [Ethernet]
    [     PW     ]
    [ -TP LSP ]

b)  [  IP  ] [  IP  ] [  IP  ]
    [ Demux  ]
    [ -TP LSP ]

Figure 1: MPLS-TP Server Layer Example
```

MPLS-TP LSP, provided by a network that also supports non-MPLS-TP functions, acting as a server for other layer 1, layer 2 and layer 3 networks.

```plaintext
<-- L1/2/3 -->|<-- MPLS -->|<-- L1/2/3 -->|

MPLS-TP
++++   LSP   ++++
++++  Client  +------ Client  ++++
|CE1=Traffic=|PE2=---------|PE3=Traffic=|CE1|
++++  +------ +------ +------ +------ +------ +++

Example
a)  [Ethernet] [Ethernet] [Ethernet]
    [     PW     ]
    [ -TP LSP ]

b)  [  IP  ] [  IP  ] [  IP  ]
    [ Demux  ]
    [ -TP LSP ]

Figure 2: MPLS-TP in MPLS Network Example
MPLS-TP as a server layer for client layer traffic of IP or MPLS networks which do not use functions of the MPLS transport profile.

MPLS-TP must not modify the MPLS forwarding architecture and must be based on existing pseudowire and LSP constructs.

Point to point LSPs may be unidirectional or bi-directional, and it must be possible to construct congruent bi-directional LSPs.

MPLS-TP LSPs do not merge with other LSPs at an MPLS-TP LSR and it must be possible to detect if a merged LSP has been created.

Figure 3: MPLS-TP Transporting Client Service Traffic

2. MPLS Transport Profile Requirements

The requirements for MPLS-TP are specified in [RFC5654], [I-D.ietf-mpls-tp-oam-requirements], and [I-D.ietf-mpls-tp-nm-req]. This section provides a brief reminder to guide the reader and is therefore not normative. It is not intended as a substitute for these documents.
It must be possible to forward packets solely based on switching the MPLS or PW label. It must also be possible to establish and maintain LSPs and/or pseudowires both in the absence or presence of a dynamic control plane. When static provisioning is used, there must be no dependency on dynamic routing or signaling.

OAM, protection and forwarding of data packets must be able to operate without IP forwarding support.

It must be possible to monitor LSPs and pseudowires through the use of OAM in the absence of control plane or routing functions. In this case information gained from the OAM functions is used to initiate path recovery actions at either the PW or LSP layers.

3. MPLS Transport Profile Overview

3.1. Packet Transport Services

One objective of MPLS-TP is to enable MPLS networks to provide packet transport services with a similar degree of predictability to that found in existing transport networks. Such packet transport services inherit a number of characteristics, defined in [RFC5654]:

- In an environment where an MPLS-TP layer network is supporting a client layer network, and the MPLS-TP layer network is supported by a server layer network then operation of the MPLS-TP layer network must be possible without any dependencies on either the server or client layer network.

- The service provided by the MPLS-TP network to the client is guaranteed not to fall below the agreed level regardless of other client activity.

- The control and management planes of any client network layer that uses the service is isolated from the control and management planes of the MPLS-TP layer network, where the client network layer is considered to be the native service of the MPLS-TP network.

- Where a client network makes use of an MPLS-TP server that provides a packet transport service, the level of co-ordination required between the client and server layer networks is minimal (preferably no co-ordination will be required).

- The complete set of packets generated by a client MPLS(-TP) layer network using the packet transport service, which may contain packets that are not MPLS packets (e.g. IP or CLNS packets used by the control/management plane of the client MPLS(-TP) layer network.
The packet transport service enables the MPLS-TP layer network addressing and other information (e.g. topology) to be hidden from any client layer networks using that service, and vice-versa.

These characteristics imply that a packet transport service does not support a connectionless packet-switched forwarding mode. However, this does not preclude it carrying client traffic associated with a connectionless service.

Such packet transport services are very similar to Layer 2 Virtual Private Networks as defined by the IETF.

### 3.2. Scope of the MPLS Transport Profile

Figure 4 illustrates the scope of MPLS-TP. MPLS-TP solutions are primarily intended for packet transport applications. MPLS-TP is a strict subset of MPLS, and comprises only those functions that are necessary to meet the requirements of [RFC5654]. This includes MPLS functions that were defined prior to [RFC5654] but that meet the requirements of [RFC5654], together with additional functions defined to meet those requirements. Some MPLS functions defined before [RFC5654] such as Equal Cost Multi-Path, LDP signaling used in such a way that it creates multipoint-to-point LSPs, and IP forwarding in the data plane are explicitly excluded from MPLS-TP by that requirements specification.

Note that MPLS as a whole will continue to evolve to include additional functions that do not conform to the MPLS Transport Profile or its requirements, and thus fall outside the scope of MPLS-TP.

```
|<================================= MPLS =============================>|

|<---------- Pre-RFC5654 MPLS ==============>|
{       ECMP       }
{     LDP/non-TE LSPs  }
{     IP fwd        }

|<================================= MPLS-TP ====================>|
{ Additional }
{   Transport }
{  Functions  }
```
3.3. Architecture

MPLS-TP comprises the following architectural elements:

- A standard MPLS data plane [RFC3031] as profiled in [I-D.fbb-mpls-tp-data-plane].
- Sections, LSPs and PWs that provide a packet transport service for a client network.
- Proactive and on-demand Operations, Administration and Maintenance (OAM) functions to monitor and diagnose the MPLS-TP network, such as connectivity check, connectivity verification, performance monitoring and fault localisation.
- Optional control planes for LSPs and PWs, as well as support for static provisioning and configuration.
- Optional path protection mechanisms to ensure that the packet transport service survives anticipated failures and degradations of the MPLS-TP network.
- Network management functions.

The MPLS-TP architecture for LSPs and PWs includes the following two sets of functions:

- MPLS-TP client adaptation
- MPLS-TP forwarding

The adaptation functions interface the native service to MPLS-TP. This includes the case where the native service is an MPLS-TP LSP.

The forwarding functions comprise the mechanisms required for forwarding the encapsulated client traffic over an MPLS-TP server layer network, for example PW and LSP labels.

3.3.1. MPLS-TP Client Adaptation Functions

The MPLS-TP native service adaptation functions interface the client service to MPLS-TP. For pseudowires, these adaptation functions are the payload encapsulation described in Section 4.4 of [RFC3985] and Section 6 of [RFC5659]. For network layer client services, the adaptation function uses the MPLS encapsulation format as defined in [RFC3032].
The purpose of this encapsulation is to abstract the client service data plane from the MPLS-TP data plane, thus contributing to the independent operation of the MPLS-TP network.

MPLS-TP is itself a client of an underlying server layer. MPLS-TP is thus also bounded by a set of adaptation functions to this server layer network, which may itself be MPLS-TP. These adaptation functions provide encapsulation of the MPLS-TP frames and for the transparent transport of those frames over the server layer network. The MPLS-TP client inherits its Quality of Service (QoS) from the MPLS-TP network, which in turn inherits its QoS from the server layer. The server layer must therefore provide the necessary QoS to ensure that the MPLS-TP client QoS commitments can be satisfied.

### 3.3.2. MPLS-TP Forwarding Functions

The forwarding functions comprise the mechanisms required for forwarding the encapsulated client over an MPLS-TP server layer network, for example PW and LSP labels.

MPLS-TP LSPs use the MPLS label switching operations and TTL processing procedures defined in [RFC3031] and [RFC3032]. These operations are highly optimised for performance and are not modified by the MPLS-TP profile.

In addition, MPLS-TP PWs use the SS-PW and MS-PW forwarding operations defined in [RFC3985] and [RFC5659]. The PW label is processed by a PW forwarder and is always at the bottom of the label stack for a given MPLS-TP layer network.

Per-platform label space is used for PWs. Either per-platform, per-interface or other context-specific label space [RFC5331] may be used for LSPs.

MPLS-TP forwarding is based on the label that identifies the transport path (LSP or PW). The label value specifies the processing operation to be performed by the next hop at that level of encapsulation. A swap of this label is an atomic operation in which the contents of the packet after the swapped label are opaque to the forwarder. The only event that interrupts a swap operation is TTL expiry. This is a fundamental architectural construct of MPLS to be taken into account when designing protocol extensions that require packets (e.g. OAM packets) to be sent to an intermediate LSR.

Further processing to determine the context of a packet occurs when a swap operation is interrupted in this manner, or a pop operation exposes a specific reserved label at the top of the stack, or the packet is received with the GAL (Section 3.6) at the top of stack.
Otherwise the packet is forwarded according to the procedures in [RFC3032].

Point-to-point MPLS-TP LSPs can be either unidirectional or bidirectional.

It must be possible to configure an MPLS-TP LSP such that the forward and backward directions of a bidirectional MPLS-TP LSP are co-routed, i.e. follow the same path. The pairing relationship between the forward and the backward directions must be known at each LSR or LER on a bidirectional LSP.

In normal conditions, all the packets sent over a PW or an LSP follow the same path through the network and those that belong to a common ordered aggregate are delivered in order. For example per-packet equal cost multi-path (ECMP) load balancing is not applicable to MPLS-TP LSPs.

Penultimate hop popping (PHP) is disabled on MPLS-TP LSPs by default.

MPLS-TP supports Quality of Service capabilities via the MPLS Differentiated Services (DiffServ) architecture [RFC3270]. Both E-LSP and L-LSP MPLS DiffServ modes are supported. The Traffic Class field (formerly the EXP field) of an MPLS label follows the definition and processing rules of [RFC5462] and [RFC3270]. Note that packet reordering between flows belonging to different traffic classes may occur if more than one traffic class is supported on a single LSP.

Only the Pipe and Short Pipe DiffServ tunnelling and TTL processing models described in [RFC3270] and [RFC3443] are supported in MPLS-TP.

3.4. MPLS-TP Native Services

This document describes the architecture for two native service adaptation mechanisms, which provide encapsulation and demultiplexing for native service traffic traversing an MPLS-TP network:

- A PW
- An MPLS Label

A PW provides any emulated service that the IETF has defined to be provided by a PW, for example Ethernet, Frame Relay, or PPP/HDLC. A registry of PW types is maintained by IANA. When the native service adaptation is via a PW, the mechanisms described in Section 3.4.2 are used.
An MPLS LSP Label can also be used as the adaptation, in which case any native service traffic type supported by [RFC3031] and [RFC3032] is allowed. Examples of such traffic types include IP, and MPLS-labeled packets. Note that the latter case includes TE-LSPs [RFC3209] and LSP based applications such as PWs, Layer 2 VPNs [RFC4664], and Layer 3 VPNs [RFC4364]. When the native service adaptation is via an MPLS label, the mechanisms described in Section 3.4.3 are used.

3.4.1. MPLS-TP Client/Server Relationship

The MPLS-TP client server relationship is defined by the MPLS-TP network boundary and the label context. It is not explicitly indicated in the packet. In terms of the MPLS label stack, when the client traffic type of the MPLS-TP network is an MPLS LSP or a PW, then the S bits of all the labels in the MPLS-TP label stack carrying that client traffic are zero; otherwise the bottom label of the MPLS-TP label stack has the S bit set to 1 (i.e. there can only one S bit set in a label stack).

The data plane behaviour of MPLS-TP is the same as the best current practise for MPLS. This includes the setting of the S-Bit. In each case, the S-bit is set to indicate the bottom (i.e. inner-most) label in the label stack that is contiguous between the MPLS-TP server and the client layer. Note that this best current practise differs slightly from [RFC3032] which uses the S-bit to identify when MPLS label processing stops and network layer processing starts.

The relationship of MPLS-TP to its clients is illustrated in Figure 5.
PW-Based Services | MPLS Labelled Services | IP Transport
----------------|----------------------|--------------
Emulated Service | PW over LSP | IP over LSP | IP

+------------+
| PW Payload |
| PW Lbl(S=1) | | IP |
+------------+
+------------+
| PW Payload |
| LSP Lbl(S=0) | | LSP Lbl(S=1) | | IP |

~~~~~~~~~~~ denotes Client <-> MPLS-TP layer boundary

Note that in the PW over LSP case the client may omit its LSP Label if penultimate hop popping has been agreed with its peer

Figure 5: MPLS-TP - Client Relationship

The data plane behaviour of MPLS-TP is the same as the best current practise for MPLS. This includes the setting of the S-Bit. In each case, the S-bit is set to indicate the bottom (i.e. inner-most) label in the label stack that is contiguous between the MPLS-TP server and the client layer.

Note that the label stacks shown above are divided between those inside the MPLS-TP Network and those within the client network when the client network is MPLS(-TP). They illustrate the smallest number of labels possible. These label stacks could also include more labels.

3.4.2. Pseudowire Adaptation

The architecture for an MPLS-TP network that provides PW emulated services is based on the MPLS [RFC3031] and pseudowire [RFC3985] architectures. Multi-segment pseudowires may optionally be used to provide a packet transport service, and their use is consistent with
the MPLS-TP architecture. The use of MS-PWs may be motivated by, for example, the requirements specified in [RFC5254]. If MS-PWs are used, then the MS-PW architecture [RFC5659] also applies.

Figure 6 shows the architecture for an MPLS-TP network using single-segment PWs.

Figure 7 shows the architecture for an MPLS-TP network when multi-segment pseudowires are used. Note that as in the SS-PW case, P-routers may also exist.
PW1-segment1 and PW1-segment2 are segments of the same MS-PW, while PW2-segment1 and PW2-segment2 are segments of another MS-PW.

Figure 7: MPLS-TP Architecture (Multi-Segment PW)

The corresponding MPLS-TP protocol stacks including PWs are shown in Figure 8. In this figure the Transport Service Layer [RFC5654] is identified by the PW demultiplexer (Demux) label and the Transport Path Layer [RFC5654] is identified by the LSP Demux Label.
3.4.2.1. Pseudowire Based Services

When providing a Virtual Private Wire Service (VPWS), Virtual Private Local Area Network Service (VPLS), Virtual Private Multicast Service (VPMS) or Internet Protocol Local Area Network Service (IPLS), pseudowires must be used to carry the client service. VPWS, VPLS, and IPLS are described in [RFC4664]. VPMS is described in [I-D.ietf-l2vpn-vpms-frmk-w-requirements].

3.4.3. Network Layer Adaptation

MPLS-TP LSPs can be used to transport network layer clients. This document uses the term Network Layer in the same sense as it is used in [RFC3031] and [RFC3032]. The network layer protocols supported by [RFC3031] and [RFC3032] can be transported between service interfaces. Examples are shown in Figure 5 above. Support for network layer clients follows the MPLS architecture for support of network layer protocols as specified in [RFC3031] and [RFC3032].

With network layer adaptation, the MPLS-TP domain provides either a uni-directional or bidirectional point-to-point connection between two PEs in order to deliver a packet transport service to attached customer edge (CE) nodes. For example, a CE may be an IP, MPLS or
MPLS-TP node. As shown in Figure 9, there is an attachment circuit between the CE node on the left and its corresponding provider edge (PE) node which provides the service interface, a bidirectional LSP across the MPLS-TP network to the corresponding PE node on the right, and an attachment circuit between that PE node and the corresponding CE node for this service.

The attachment circuits may be heterogeneous (e.g., any combination of SDH, PPP, Frame Relay, etc.) and network layer protocol payloads arrive at the service interface encapsulated in the Layer1/Layer2 encoding defined for that access link type. It should be noted that the set of network layer protocols includes MPLS and hence MPLS encoded packets with an MPLS label stack (the client MPLS stack), may appear at the service interface.

<table>
<thead>
<tr>
<th>Client Network Layer</th>
</tr>
</thead>
<tbody>
<tr>
<td>Pkt Xport Service</td>
</tr>
<tr>
<td>PSN Tunnel</td>
</tr>
<tr>
<td>V V V V V V</td>
</tr>
<tr>
<td>AC PE1 LSP CE1 IP</td>
</tr>
<tr>
<td>AC PE2 LSP CE2 IP</td>
</tr>
<tr>
<td>^ Provider Edge</td>
</tr>
<tr>
<td>^ Provider Edge</td>
</tr>
<tr>
<td>Customer 1 P-router</td>
</tr>
<tr>
<td>Edge 1</td>
</tr>
<tr>
<td>LSP TE</td>
</tr>
<tr>
<td>Edge 2</td>
</tr>
<tr>
<td>LSP TE</td>
</tr>
<tr>
<td>Customer 2</td>
</tr>
</tbody>
</table>

Figure 9: MPLS-TP Architecture for Network Layer Clients

At the ingress service interface the client packets are received. The PE pushes one or more labels onto the client packets which are then label switched over the transport network. Correspondingly the egress PE pops any labels added by the MPLS-TP networks and transmits the packet for delivery to the attached CE via the egress service interface.
Figure 10: Domain of MPLS-TP Layer Network for IP and LSP Clients

In this figure the Transport Service Layer [RFC5654] is identified by the Service LSP (SvcLSP) demultiplexer (Demux) label and the Transport Path Layer [RFC5654] is identified by the LSP Demux Label. Note that the functions of the Encapsulation label and the Service Label shown above as SvcLSP Demux may be represented by a single label stack entry. Additionally, the S-bit will always be zero when the client layer is MPLS labelled.

Within the MPLS-TP transport network, the network layer protocols are carried over the MPLS-TP network using a logically separate MPLS label stack (the server stack). The server stack is entirely under the control of the nodes within the MPLS-TP transport network and it is not visible outside that network. Figure 10 shows how a client network protocol stack (which may be an MPLS label stack and payload) is carried over a network layer client service over an MPLS-TP transport network.

A label per network layer protocol payload type that is to be transported is required. When multiple protocol payload types are to be carried over a single service a unique label stack entry must be present for each payload type. Such labels are referred to as "Encapsulation Labels", one of which is shown in Figure 10. Encapsulation Label may be either configured or signaled.

Both an Encapsulation Label and a Service Label should be present in
the label stack when a particular packet transport service is supporting more than one network layer protocol payload type. For example, if both IP and MPLS are to be carried, as shown in Figure 9, then two Encapsulation Labels are mapped on to a common Service Label.

Note: The Encapsulation Label may be omitted when the transport service is supporting only one network layer protocol payload type. For example, if only MPLS labeled packets are carried over a service, then the Service Label (stack entry) provides both the payload type indication and service identification.

Service labels are typically carried over an MPLS-TP LSP edge-to-edge (or transport path layer). An MPLS-TP edge-to-edge LSP is represented as an LSP Demux label as shown in Figure 10. An edge-to-edge LSP is commonly used when more than one service exists between two PEs.

Note that the edge-to-edge LSP may be omitted when only one service exists between two PEs. For example, if only one service is carried between two PEs then a single Service Label could be used to provide both the service indication and the MPLS-TP edge-to-edge LSP. Alternatively, if multiple services exist between a pair of PEs then a per-client Service Label would be mapped on to a common MPLS-TP edge-to-edge LSP.

As noted above, the layer 2 and layer 1 protocols used to carry the network layer protocol over the attachment circuits are not transported across the MPLS-TP network. This enables the use of different layer 2 and layer 1 protocols on the two attachment circuits.

At each service interface, Layer 2 addressing must be used to ensure the proper delivery of a network layer packet to the adjacent node. This is typically only an issue for LAN media technologies (e.g., Ethernet) which have Media Access Control (MAC) addresses. In cases where a MAC address is needed, the sending node must set the destination MAC address to an address that ensures delivery to the adjacent node. That is the CE sets the destination MAC address to an address that ensures delivery to the PE, and the PE sets the destination MAC address to an address that ensures delivery to the CE. The specific address used is technology type specific and is not specified in this document. In some technologies the MAC address will need to be configured. (Examples for the Ethernet case include a configured unicast MAC address for the adjacent node, or even using the broadcast MAC address when the CE-PE service interface is dedicated. The configured address is then used as the destination MAC address for all packets sent over the service interface.)
Note that when two CEs, which peer with each other, operate over a network layer transport service and run a routing protocol such as IS-IS or OSPF, some care should be taken to configure the routing protocols to use point-to-point adjacencies. The specifics of such configuration is outside the scope of this document. See [RFC5309] for additional details.

The CE to CE service types and corresponding labels may be configured or signaled. See Section 3.11 for additional details related to configured service types. See Section 3.9 for additional details related to signaled service types.

3.5. Identifiers

Identifiers are used to uniquely distinguish entities in an MPLS-TP network. These include operators, nodes, LSPs, pseudowires, and their associated maintenance entities. [I-D.ietf-mpls-tp-identifiers] defines a set of identifiers that are compatible with existing MPLS control plane identifiers, as well as a set of identifiers that may be used when no IP control plane is available.

3.6. Generic Associated Channel (G-ACh)

For correct operation of the OAM it is important that the OAM packets fate-share with the data packets. In addition in MPLS-TP it is necessary to discriminate between user data payloads and other types of payload. For example, a packet may be associated with a Signaling Communication Channel (SCC), or a channel used for Automatic Protection Switching (APS) data. This is achieved by carrying such packets on a generic control channel associated to the LSP, PW or section.

MPLS-TP makes use of such a generic associated channel (G-ACh) to support Fault, Configuration, Accounting, Performance and Security (FCAPS) functions by carrying packets related to OAM, APS, SCC, MCC or other packet types in-band over LSPs or PWs. The G-ACh is defined in [RFC5586] and is similar to the Pseudowire Associated Channel [RFC4385], which is used to carry OAM packets over pseudowires. The G-ACh is indicated by a generic associated channel header (ACH), similar to the Pseudowire VCCV control word; this header is present for all Sections, LSPs and PWs making use of FCAPS functions supported by the G-ACh.

For pseudowires, the G-ACh uses the first four bits of the pseudowire control word to provide the initial discrimination between data packets and packets belonging to the associated channel, as described in [RFC4385]. When this first nibble of a packet, immediately
following the label at the bottom of stack, has a value of ‘1’, then
this packet belongs to a G-ACh. The first 32 bits following the
bottom of stack label then have a defined format called an associated
channel header (ACH), which further defines the content of the
packet. The ACH is therefore both a demultiplexer for G-ACh traffic
on the PW, and a discriminator for the type of G-ACh traffic.

When the OAM or other control message is carried over an LSP, rather
than over a pseudowire, it is necessary to provide an indication in
the packet that the payload is something other than a user data
packet. This is achieved by including a reserved label with a value
of 13 in the label stack. This reserved label is referred to as the
‘G-ACh Label (GAL)’, and is defined in [RFC5586]. When a GAL is
found, it indicates that the payload begins with an ACH. The GAL is
thus a demultiplexer for G-ACh traffic on the LSP, and the ACH is a
discriminator for the type of traffic carried on the G-ACh. Note
however that MPLS-TP forwarding follows the normal MPLS model, and
that a GAL is invisible to an LSR unless it is the top label in the
label stack. The only other circumstance under which the label stack
may be inspected for a GAL is when the TTL has expired. Any MPLS-TP
component that intentionally performs this inspection must assume
that it is asynchronous with respect to the forwarding of other
packets. All operations on the label stack are in accordance with
[RFC3031] and [RFC3032].

In MPLS-TP, the ‘G-ACh Label (GAL)’ always appears at the bottom of
the label stack (i.e. its S bit is set to 1).

The G-ACh must only be used for channels that are an adjunct to the
data service. Examples of these are OAM, APS, MCC and SCC, but the
use is not restricted to these services. The G-ACh must not be used
to carry additional data for use in the forwarding path, i.e. it must
not be used as an alternative to a PW control word, or to define a PW
type.

At the server layer, bandwidth and QoS commitments apply to the gross
traffic on the LSP, PW or section. Since the G-ACh traffic is
indistinguishable from the user data traffic, protocols using the
G-ACh must take into consideration the impact they have on the user
data that they are sharing resources with. Conversely, capacity must
be made available for important G-ACh uses such as protection and
OAM. In addition, protocols using the G-ACh must conform to the
security and congestion considerations described in [RFC5586].

Figure 11 shows the reference model depicting how the control channel
is associated with the pseudowire protocol stack. This is based on
the reference model for VCCV shown in Figure 2 of [RFC5085].
Figure 11: PWE3 Protocol Stack Reference Model showing the G-ACh PW associated channel messages are encapsulated using the PWE3 encapsulation, so that they are handled and processed in the same manner (or in some cases, an analogous manner) as the PW PDUs for which they provide a control channel.

Figure 12 shows the reference model depicting how the control channel is associated with the LSP protocol stack.
3.7. Operations, Administration and Maintenance (OAM)

MPLS-TP must be able to operate in environments where IP is not used in the forwarding plane. Therefore, the default mechanism for OAM demultiplexing in MPLS-TP LSPs and PWs is the Generic Associated Channel (Section 3.6). Forwarding based on IP addresses for user or OAM packets is not required for MPLS-TP.

[RFC4379] and BFD for MPLS LSPs [I-D.ietf-bfd-mpls] have defined alert mechanisms that enable an MPLS LSR to identify and process MPLS OAM packets when the OAM packets are encapsulated in an IP header. These alert mechanisms are based on TTL expiration and/or use an IP destination address in the range 127/8 for IPv4 and that same range embedded as IPv4 mapped IPv6 addresses for IPv6 [RFC4379]. When the OAM packets are encapsulated in an IP header, these mechanisms are the default mechanisms for MPLS networks in general for identifying MPLS OAM packets. MPLS-TP must be able to operate in an environments where IP forwarding is not supported, and thus the G-ACh/GAL is the default mechanism to demultiplex OAM packets in MPLS-TP.

MPLS-TP supports a comprehensive set of OAM capabilities for packet transport applications, with equivalent capabilities to those
provided in SONET/SDH.

MPLS-TP defines mechanisms to differentiate specific packets (e.g. OAM, APS, MCC or SCC) from those carrying user data packets on the same transport path (i.e. section, LSP or PW). These mechanisms are described in [RFC5586].

MPLS-TP requires [I-D.ietf-mpls-tp-oam-requirements] that a set of OAM capabilities is available to perform fault management (e.g. fault detection and localisation) and performance monitoring (e.g. packet delay and loss measurement) of the LSP, PW or section. The framework for OAM in MPLS-TP is specified in [I-D.ietf-mpls-tp-oam-framework].

MPLS-TP OAM packets share the same fate as their corresponding data packets, and are identified through the Generic Associated Channel mechanism [RFC5586]. This uses a combination of an Associated Channel Header (ACH) and a G-ACh Label (GAL) to create a control channel associated to an LSP, Section or PW.

OAM and monitoring in MPLS-TP is based on the concept of maintenance entities, as described in [I-D.ietf-mpls-tp-oam-framework]. A Maintenance Entity can be viewed as the association of two Maintenance End Points (MEPs). A Maintenance Entity Group (MEG) is a collection of one or more MEs that belongs to the same transport path and that are maintained and monitored as a group. The MEPs that form an ME limit the OAM responsibilities of an OAM flow to within the domain of a transport path or segment, in the specific layer network that is being monitored and managed.

An ME may also include a set of Maintenance Intermediate Points (MIPs). Maintenance End Points (MEPs) are capable of sourcing and sinking OAM flows, while Maintenance Intermediate Points (MIPs) can only sink or respond to OAM flows from within a MEG, or originate notifications as a result of specific network conditions.

The following MPLS-TP MEs are specified in [I-D.ietf-mpls-tp-oam-framework]:

- A Section Maintenance Entity (SME), allowing monitoring and management of MPLS-TP Sections (between MPLS LSRs).
- A LSP Maintenance Entity (LME), allowing monitoring and management of an edge-to-edge LSP (between LERs).
- A PW Maintenance Entity (PME), allowing monitoring and management of an edge-to-edge SS/MS-PWs (between T-PEs).
A G-ACh packet may be directed to an individual MIP along the path of an LSP or MS-PW by setting the appropriate TTL in the label for the G-ACh packet, as per the traceroute mode of LSP Ping [RFC4379] and the vccv-trace mode of [I-D.ietf-pwe3-segmented-pw]. Note that this works when the location of MIPs along the LSP or PW path is known by the MEP. There may be circumstances where this is not the case, e.g. following restoration using a facility bypass LSP. In these cases, tools to trace the path of the LSP may be used to determine the appropriate setting for the TTL to reach a specific MIP.

Within an LSR or PE, MEPs and MIPs can only be placed where MPLS layer processing is performed on a packet. The architecture mandates that this must occur at least once.

MEPs may only act as a sink of OAM packets when the label associated with the LSP or PW for that ME is popped. MIPs can only be placed where an exception to the normal forwarding operation occurs. A MEP may act as a source of OAM packets wherever a label is pushed or swapped. For example, on an MS-PW, a MEP may source OAM within an S-PE or a T-PE, but a MIP may only be associated with a S-PE and a sink MEP can only be associated with a T-PE.

The MPLS-TP OAM architecture supports a wide range of OAM functions to check continuity, to verify connectivity and to monitor the performance of the path, to generate, filter and manage local and remote defect alarms. These functions are applicable to any layer defined within MPLS-TP, i.e. to MPLS-TP Sections, LSPs and PWs.

The MPLS-TP OAM tool-set must be able to operate without relying on a dynamic control plane or IP functionality in the datapath. In the case of an MPLS-TP deployment in a network in which IP functionality is available, all existing IP/MPLS OAM functions, e.g. LSP-Ping, BFD and VCCV, may be used.

3.8. LSP Return Path

Management, control and OAM protocol functions may require response packets to be delivered from the receiver back to the originator of a message exchange. This section provides a summary of the return path options in MPLS-TP networks.

In this discussion we assume that A and B are terminal LSRs (i.e. LERs) for an MPLS-TP LSP and that Y is an intermediate LSR along the LSP. In the unidirectional case, A is taken to be the upstream and B the downstream LSR with respect to the LSP. We consider the following cases for the various types of LSPs:
1. Packet transmission from B to A
2. Packet transmission from Y to A
3. Packet transmission from B to Y

Note that a return path may not always exist, and that packet transmission in one or more of the above cases may not be possible. In general the existence and nature of return paths for MPLS-TP LSPs is determined by operational provisioning.

3.8.1. Return Path Types

There are two types of return path that may be used for the delivery of traffic from a downstream node D to an upstream node U either:

a. D maintains an MPLS-TP LSP back to U which is specifically designated to carry return traffic for the original LSP, or

b. D has some other unspecified means of directing traffic back to U.

The first option is referred to as an "in-band" return path, the second as an "out-of-band" return path.

There are various possibilities for "out-of-band" return paths. Such a path may, for example, be based on ordinary IP routing. In this case packets would be forwarded as usual to a destination IP address associated with U. In an MPLS-TP network that is also an IP/MPLS network, such a forwarding path may traverse the same physical links or logical transport paths used by MPLS-TP. An out-of-band return path may also be indirect, via a distinct Data Communication Network (DCN) (provided, for example, by the method specified in [RFC5718]); or it may be via one or more other MPLS-TP LSPs.

3.8.2. Point-to-Point Unidirectional LSPs

Case 1 In this situation, either an in-band or out-of-band return path may be used to deliver traffic from B back to A.

In the in-band case there is in essence an associated bidirectional LSP between A and B, and the discussion for such LSPs below applies. It is therefore recommended for reasons of operational simplicity that point-to-point unidirectional LSPs be provisioned as associated bidirectional LSPs (which may also be co-routed) whenever return traffic from B to A is required. Note that the two directions of such an LSP may have differing bandwidth
allocations and QoS characteristics.

Case 2  In this case only the out-of-band return path option is available. However, an additional out-of-band possibility is worthy of note here: if B is known to have a return path to A, then Y can arrange to deliver return traffic to A by first sending it to B along the original LSP. The mechanism by which B recognises the need for and performs this forwarding operation is protocol-specific.

Case 3  In this case only the out-of-band return path option is available. However, if B has a return path to A, then in a manner analogous to the previous case B can arrange to deliver return traffic to Y by first sending it to A along that return path. The mechanism by which A recognises the need for and performs this forwarding operation is protocol-specific.

3.8.3. Point-to-Point Associated Bidirectional LSPs

For Case 1, B has a natural in-band return path to A, the use of which is typically preferred for return traffic, although out-of-band return paths are also applicable.

For Cases 2 and 3, the considerations are the same as those for point-to-point unidirectional LSPs.

3.8.4. Point-to-Point Co-Routed Bidirectional LSPs

For all of Cases 1, 2, and 3, a natural in-band return path exists in the form of the LSP itself, and its use is typically preferred for return traffic. Out-of-band return paths, however, are also applicable, primarily as an alternative means of delivery in case the in-band return path has failed.

3.9. Control Plane

A distributed dynamic control plane may be used to enable dynamic service provisioning in an MPLS-TP network. Where the requirements specified in [RFC5654] can be met, the MPLS Transport Profile uses existing standard control plane protocols for LSPs and PWs.

Note that a dynamic control plane is not required in an MPLS-TP network. See Section 3.11 for further details on statically configured and provisioned MPLS-TP services.

Figure 13 illustrates the relationship between the MPLS-TP control plane, the forwarding plane, the management plane, and OAM for point-
to-point MPLS-TP LSPs or PWs.

+------------------------------------------------------------------+
|                                                                  |
|                Network Management System and/or                   |
| Control Plane for Point to Point Connections                    |
+------------------------------------------------------------------+

|          |          |          |          |          |          |          |          |          |          |

Note:
1) NMS may be centralised or distributed. Control plane is distributed.
2) ‘Edge’ functions refers to those functions present at the edge of a PSN domain, e.g. NSP or classification.
3) The control plane may be transported over the server layer, an LSP or a G-ACh.

Figure 13: MPLS-TP Control Plane Architecture Context

The MPLS-TP control plane is based on existing MPLS and PW control plane protocols. MPLS-TP uses Generalized MPLS (GMPLS) signaling ([RFC3945], [RFC3471], [RFC3473]) for LSPs and Targeted LDP (T-LDP) [RFC4447] [I-D.ietf-pwe3-segmented-pw][I-D.ietf-pwe3-dynamic-ms-pw] for pseudowires.

MPLS-TP requires that any signaling be capable of being carried over an out-of-band signaling network or a signaling control channel such as the one described in [RFC5718]. Note that while T-LDP signaling is traditionally carried in-band in IP/MPLS networks, this does not preclude its operation over out-of-band channels. References to T-LDP in this document do not preclude the definition of alternative PW control protocols for use in MPLS-TP.

PW control (and maintenance) takes place separately from LSP tunnel signaling. The main coordination between LSP and PW control will
occur within the nodes that terminate PWs. The control planes for
PWs and LSPs may be used independently, and one may be employed
without the other. This translates into the four possible scenarios:
(1) no control plane is employed; (2) a control plane is used for
both LSPs and PWs; (3) a control plane is used for LSPs, but not PWs;
(4) a control plane is used for PWs, but not LSPs. The PW and LSP
control planes, collectively, must satisfy the MPLS-TP control plane
requirements reviewed in the MPLS-TP Control Plane Framework
[I-D.abfb-mpls-tp-control-plane-framework]. When client services
are provided directly via LSPs, all requirements must be satisfied by the
LSP control plane. When client services are provided via PWs, the PW
and LSP control planes operate in combination and some functions may
be satisfied via the PW control plane while others are provided to
PWs by the LSP control plane.

Note that if MPLS-TP is being used in a multi-layer network, a number
of control protocol types and instances may be used. This is
consistent with the MPLS architecture which permits each label in the
label stack to be allocated and signaled by its own control protocol.

The distributed MPLS-TP control plane may provide the following
functions:

- Signaling
- Routing
- Traffic engineering and constraint-based path computation

In a multi-domain environment, the MPLS-TP control plane supports
different types of interfaces at domain boundaries or within the
domains. These include the User-Network Interface (UNI), Internal
Network Node Interface (I-NNI), and External Network Node Interface
(E-NNI). Note that different policies may be defined that control
the information exchanged across these interface types.

The MPLS-TP control plane is capable of activating MPLS-TP OAM
functions as described in the OAM section of this document
Section 3.7, e.g. for fault detection and localisation in the event
of a failure in order to efficiently restore failed transport paths.

The MPLS-TP control plane supports all MPLS-TP data plane
connectivity patterns that are needed for establishing transport
paths, including protected paths as described in Section 3.12.
Examples of the MPLS-TP data plane connectivity patterns are LSPs
utilising the fast reroute backup methods as defined in [RFC4090] and
ingress-to-egress 1+1 or 1:1 protected LSPs.
The MPLS-TP control plane provides functions to ensure its own survivability and to enable it to recover gracefully from failures and degradations. These include graceful restart and hot redundant configurations. Depending on how the control plane is transported, varying degrees of decoupling between the control plane and data plane may be achieved.

3.10. Inter-domain Connectivity

A number of methods exist to support inter-domain operation of MPLS-TP, for example:

- Inter-domain TE LSPs [RFC4216]
- Multi-segment Pseudowires [RFC5659]
- LSP stitching [RFC5150]
- back-to-back attachment circuits [RFC5659]

An important consideration in selecting an inter-domain connectivity mechanism is the degree of layer network isolation and types of OAM required by the operator. The selection of which technique to use in a particular deployment scenario is outside the scope of this document.

3.11. Static Operation of LSPs and PWs

A PW or LSP may be statically configured without the support of a dynamic control plane. This may be either by direct configuration of the PEs/LSRs, or via a network management system. Static operation is independent for a specific PW or LSP instance. Thus it should be possible for a PW to be statically configured, while the LSP supporting it is set up by a dynamic control plane. When static configuration mechanisms are used, care must be taken to ensure that loops are not created.

3.12. Survivability

Survivability requirements for MPLS-TP are specified in [I-D.ietf-mpls-tp-survive-fwk].

A wide variety of resiliency schemes have been developed to meet the various network and service survivability objectives. For example, as part of the MPLS/PW paradigms, MPLS provides methods for local repair using back-up LSP tunnels ([RFC4090]), while pseudowire redundancy [I-D.ietf-pwe3-redundancy] supports scenarios where the protection for the PW cannot be fully provided by the underlying LSP.
(i.e. where the backup PW terminates on a different target PE node than the working PW in dual homing scenarios, or where protection of the S-PE is required). Additionally, GMPLS provides a well-known set of control plane driven protection and restoration mechanisms [RFC4872]. MPLS-TP provides additional protection mechanisms that are optimised for both linear topologies and ring topologies, and that operate in the absence of a dynamic control plane. These are specified in [I-D.ietf-mpls-tp-survive-fwk].

Different protection schemes apply to different deployment topologies and operational considerations. Such protection schemes may provide different levels of resiliency, for example:

- Two concurrent traffic paths (1+1).
- One active and one standby path with guaranteed bandwidth on both paths (1:1).
- One active path and a standby path the resources or which are shared by one or more other active paths (shared protection).

The applicability of any given scheme to meet specific requirements is outside the current scope of this document.

The characteristics of MPLS-TP resiliency mechanisms are as follows:

- Optimised for linear, ring or meshed topologies.
- Use OAM mechanisms to detect and localise network faults or service degenerations.
- Include protection mechanisms to coordinate and trigger protection switching actions in the absence of a dynamic control plane. This is known as an Automatic Protection Switching (APS) mechanism.
- MPLS-TP recovery schemes are applicable to all levels in the MPLS-TP domain (i.e. MPLS section, LSP and PW), providing segment and end-to-end recovery.
- MPLS-TP recovery mechanisms support the coordination of protection switching at multiple levels to prevent race conditions occurring between a client and its server layer.
- MPLS-TP recovery mechanisms can be data plane, control plane or management plane based.
- MPLS-TP supports revertive and non-revertive behaviour.
3.13. Path Segment Tunnels

In order to monitor, protect and manage a portion of an LSP, a new architectural element is defined called the Path Segment Tunnel (PST). A PST is a hierarchical LSP [RFC3031] which is defined and used for the purposes of OAM monitoring, protection or management of LSP segments or concatenated LSP segments.

A PST is defined between the edges of the portion of the LSP that needs to be monitored, protected or managed. Maintenance messages can be initiated at the edge of the PST and sent to the peer edge of the PST or to an intermediate point along the PST by setting the TTL value at the PST level accordingly.

For example in Figure 14, three PSTs are configured to allow monitoring, protection and management of the LSP concatenated segments. One PST is defined between PE1 and PE2, the second between PE2 and PE3 and a third PST is set up between PE3 and PE4. Each of these three PSTs may be monitored, protected, or managed independently.

The end-to-end traffic of the LSP, including data traffic and control traffic (OAM, Protection Switching Control, management and signaling messages) is tunneled within the PST by means of label stacking as defined in [RFC3031].

The mapping between an LSP and a PST can be 1:1, in which case it is similar to the ITU-T Tandem Connection element [G.805]. The mapping can also be 1:N to allow aggregated monitoring, protection and management of a set of LSP segments or concatenated LSP segments. Figure 15 shows a PST which is used to aggregate a set of concatenated LSP segments for the LSP from PEx to PEt and the LSP from PEa to PEd. Note that such a construct is useful, for example, when the LSPs traverse a common portion of the network and they have the same Traffic Class.
3.13.1. Provisioning of PST

PSTs can be provisioned either statically or using control plane signaling procedures. The make-before-break procedures which are supported by MPLS allow the creation of a PST on existing LSPs in-service without traffic disruption. A PST can be defined corresponding to one or more end-to-end tunneled LSPs. New end-to-end LSPs which are tunneled within the PST can be set up. Traffic of the existing LSPs is switched over to the new end-to-end tunneled LSPs. The old end-to-end LSPs can then be torn down.

3.14. Pseudowire Segment Tunnels

Pseudowire segment tunnels are for further study.

3.15. Network Management

The network management architecture and requirements for MPLS-TP are specified in [I-D.ietf-mpls-tp-nm-framework] and [I-D.ietf-mpls-tp-nm-req]. These derive from the generic specifications described in ITU-T G.7710/Y.1701 [G.7710] for transport technologies. It also incorporates the OAM requirements for MPLS Networks [RFC4377] and MPLS-TP Networks [I-D.ietf-mpls-tp-oam-requirements] and expands on those requirements to cover the modifications necessary for fault, configuration, performance, and security in a transport network.

The Equipment Management Function (EMF) of an MPLS-TP Network Element (NE) (i.e. LSR, LER, PE, S-PE or T-PE) provides the means through which a management system manages the NE. The Management Communication Channel (MCC), realised by the G-ACh, provides a logical operations channel between NEs for transferring Management information. For the management interface from a management system to an MPLS-TP NE, there is no restriction on which management
protocol is used. The MCC is used to provision and manage an end-to-end connection across a network where some segments are created/managed by, for example, Netconf [RFC4741] or SNMP [RFC3411] and other segments by XML or CORBA interfaces. Maintenance operations are run on a connection (LSP or PW) in a manner that is independent of the provisioning mechanism. An MPLS-TP NE is not required to offer more than one standard management interface. In MPLS-TP, the EMF must be capable of statically provisioning LSPs for an LSR or LER, and PWs for a PE, as well as any associated MEPs and MIPs, as per Section 3.11.

Fault Management (FM) functions within the EMF of an MPLS-TP NE enable the supervision, detection, validation, isolation, correction, and alarm handling of abnormal conditions in the MPLS-TP network and its environment. FM must provide for the supervision of transmission (such as continuity, connectivity, etc.), software processing, hardware, and environment. Alarm handling includes alarm severity assignment, alarm suppression/aggregation/correlation, alarm reporting control, and alarm reporting.

Configuration Management (CM) provides functions to control, identify, collect data from, and provide data to MPLS-TP NEs. In addition to general configuration for hardware, software protection switching, alarm reporting control, and date/time setting, the EMF of the MPLS-TP NE also supports the configuration of maintenance entity identifiers (such as MEP ID and MIP ID). The EMF also supports the configuration of OAM parameters as a part of connectivity management to meet specific operational requirements. These may specify whether the operational mode is one-time on-demand or is periodic at a specified frequency.

The Performance Management (PM) functions within the EMF of an MPLS-TP NE support the evaluation and reporting of the behaviour of the NEs and the network. One particular requirement for PM is to provide coherent and consistent interpretation of the network behaviour in a hybrid network that uses multiple transport technologies. Packet loss measurement and delay measurements may be collected and used to detect performance degradation. This is reported via fault management to enable corrective actions to be taken (e.g. protection switching), and via performance monitoring for Service Level Agreement (SLA) verification and billing. Collection mechanisms for performance data should be capable of operating on-demand or pro-actively.

4. Security Considerations

The introduction of MPLS-TP into transport networks means that the security considerations applicable to both MPLS and PWE3 apply to
those transport networks. Furthermore, when general MPLS networks that utilise functionality outside of the strict MPLS Transport Profile are used to support packet transport services, the security considerations of that additional functionality also apply.

For pseudowires, the security considerations of [RFC3985] and [RFC5659] apply.

Packets that arrive on an interface with a given label value should not be forwarded unless that label value is assigned to an LSP or PW to a peer LSR or PE that is reachable via that interface.

Each MPLS-TP solution must specify the additional security considerations that apply. This is discussed further in [I-D.fang-mpls-tp-security-framework].

5. IANA Considerations

IANA considerations resulting from specific elements of MPLS-TP functionality will be detailed in the documents specifying that functionality.

This document introduces no additional IANA considerations in itself.

6. Acknowledgements

The editors wish to thank the following for their contribution to this document:

 o Rahul Aggarwal
 o Dieter Beller
 o Malcolm Betts
 o Italo Busi
 o John E Drake
 o Hing-Kam Lam
 o Marc Lasserre
 o Vincenzo Sestito
 o Nurit Sprecher
7. Open Issues

This section contains a list of issues that must be resolved before last call.

8. References

8.1. Normative References


8.2. Informative References


Muley, P. and V. Place, "Pseudowire (PW) Redundancy", draft-ietf-pwe3-redundancy-02 (work in progress),

Muley, P. and V. Place, "Pseudowire (PW) Redundancy", draft-ietf-pwe3-redundancy-02 (work in progress),


Authors’ Addresses

Matthew Bocci (editor)
Alcatel-Lucent
Voyager Place, Shoppenhangers Road
Maidenhead, Berks SL6 2PJ
United Kingdom

Phone:
EMail: matthew.bocci@alcatel-lucent.com